

Final Exam – Solutions

Question 1:

If the type checker tries to check the statement `C().m(5,6)` then it will have to consult the symbol table of class `C` to get the type of the method `m`.

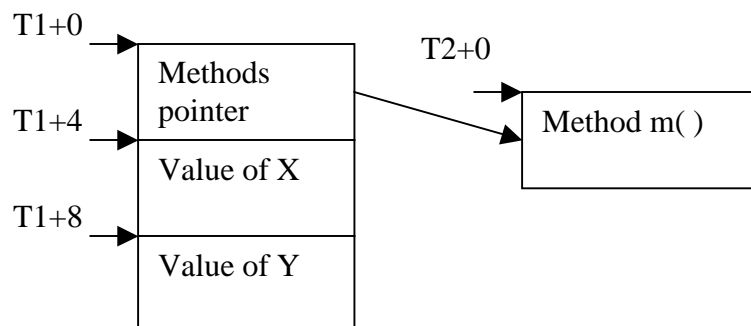
Assuming that we have a 2 pass compiler, the symbol table of class `C` should at this time look like:

Key	Value
<code>C.x</code>	<code>Int</code>
<code>C.y</code>	<code>Int</code>
<code>C.m.num</code>	<code>Int</code>
<code>C.m.a</code>	<code>Int</code>
<code>C.m.i</code>	<code>Int</code>
<code>C.m.j</code>	<code>Int</code>
Methods	
<code>C.m</code>	<code>Int</code>

Then the type check would just require access to the method field `C.m`

Question 2:

Heap layout of the `C` object where `T1` and `T2` are addresses allocated by `HALLOCATE`



Piglet translation of the code:

PRINT CALL

BEGIN

MOVE TEMP 24

BEGIN

MOVE TEMP 25 HALLOCATE 4

MOVE TEMP 26 HALLOCATE 16

HSTORE TEMP 25 0 C_m

MOVE TEMP 27 4

L0 CJUMP LT TEMP 27 16 L1

HSTORE PLUS TEMP 26 TEMP 27 0 0

MOVE TEMP 27 PLUS TEMP 27 4

JUMP L0

L1 HSTORE TEMP 26 0 TEMP 25

RETURN

TEMP 26

END

HLOAD TEMP 22 TEMP 24 0

MOVE TEMP 23 PLUS TEMP 22 0

RETURN

TEMP 23

END

(TEMP 24 5 6)

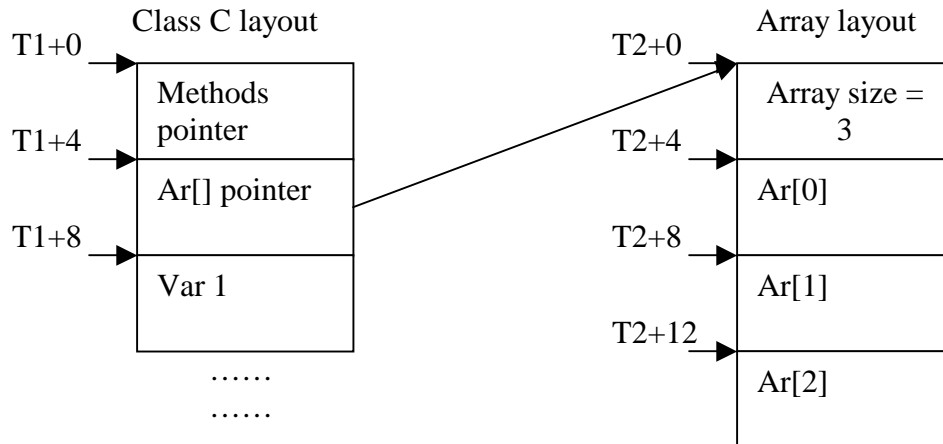
END

Question 3:

Let `int[] ar` be the first field declaration in a class `C`. Then the offset of the pointer to the array structure will be 4 (from the beginning of the class structure)

In other words if the class `C` heap layout starts at `TEMP 0` then the array pointer will be at `TEMP 0 + 4`

Assuming `HALLOCATE` returned address `T2` we do the following allocation for the array:



If we want to access the element `a[3]` and suppose that the pointer to the Array layout is stored in class variable `TEMP 20` then the Piglet code for that would look like:

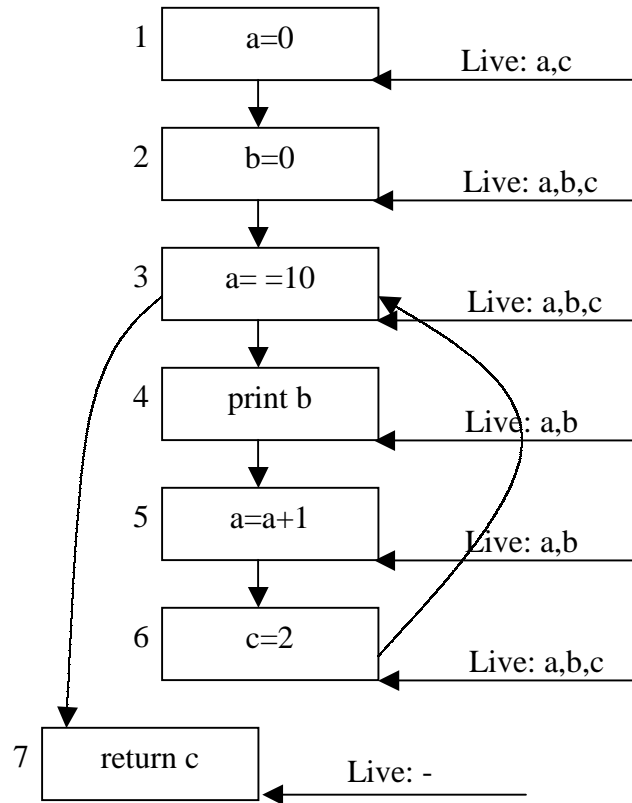
```

MOVE TEMP 21
  BEGIN
    HLOAD TEMP 32 PLUS TEMP 20 PLUS
      BEGIN
        MOVE TEMP 30 TIMES 3 4
        HLOAD TEMP 31 TEMP 20 0
        CJUMP MINUS 1 LT TEMP 30 TEMP 31 L4
        ERROR
        L4 NOOP
      RETURN
      TEMP 30
    END
    4 0
  RETURN
  TEMP 32
END

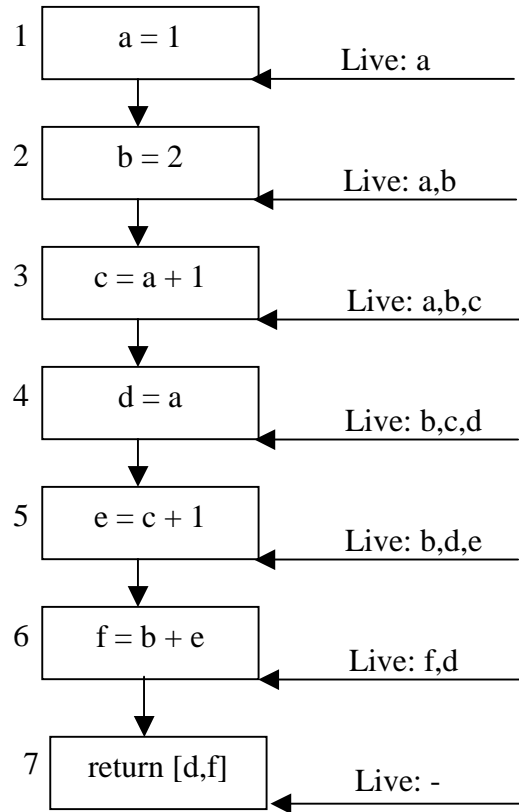
RETURN
TEMP 21

```


Then we complete the diagram:



Question 5



As we see from the diagram above we can always have at most 3 live variables which we can store in the three registers without spilling.

The following coloring scheme simplifies the initial program to:

```

r1=1
r3=2
r2=r1+1
r1=r1
r2=r2+1
r3=r3+r2
return r1,r3
    
```

